Wireless Sensor Networks 18th Lecture 10.01.2007



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Topology Control

> Networks can be too dense – too many nodes in close (radio) vicinity

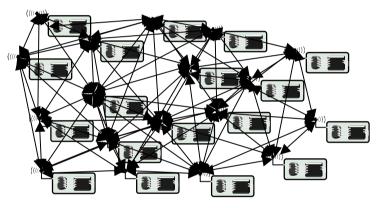
This chapter looks at methods to deal with such networks by

- Reducing/controlling transmission power
- Deciding which links to use
- Turning some nodes off



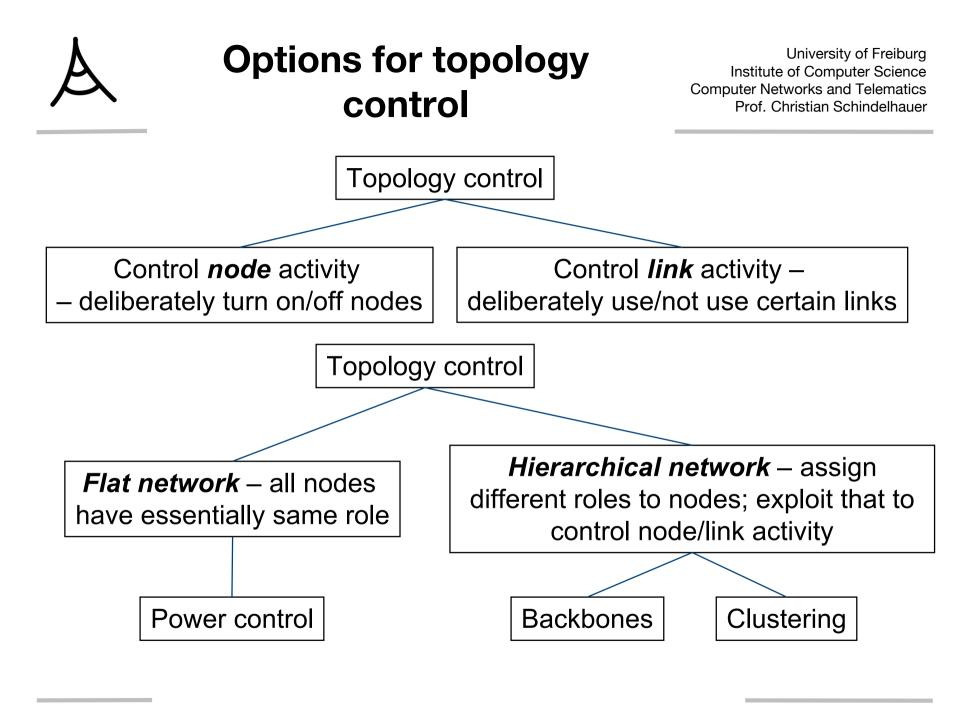
In a very dense networks, too many nodes might be in range for an efficient operation

 Too many collisions/too complex operation for a MAC protocol, too many paths to chose from for a routing protocol, …



Idea: Make topology less complex

- *Topology*: Which node is able/allowed to communicate with which other nodes
- Topology control needs to maintain invariants, e.g., connectivity



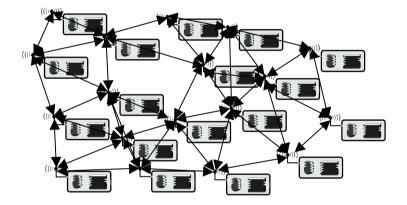


Flat networks

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Main option: Control transmission power

- Do not always use maximum power
- Selectively for some links or for a node as a whole
- Topology looks "thinner"
- Less interference, ...



Alternative: Selectively discard some links

- Usually done by introducing hierarchies



Geometric Spanners with Applications in Wireless Networks

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- 1. Introduction
 - Definition of Geometric Spanners
 - Motivation
 - Related Work
- 2. Spanners versus Weak Spanners
- 3. Spanners versus Power Spanners
- 4. Weak Spanners versus Power Spanners
 - Weak Spanners are Power Spanners if
 - Weak Spanners are Power Spanners if
 - Weak Spanners are not always Power Spanners if
 - Fractal Dimensions
- 5. Applications in Wireless Networks
- 6. Conclusions

Exponent > Dimension

Exponent = Dimension

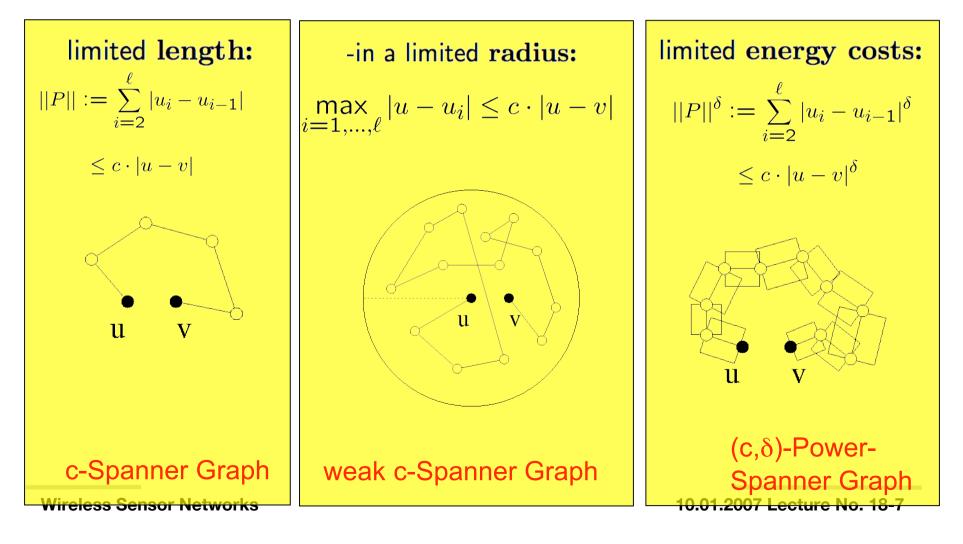
Exponent < Dimension



Geometric Spanner Graphs

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A Graph G = (V, E) with $V \subseteq \mathbf{R}$ where for all $u, v \in V$ there exists a path $P = (u = u_1, u_2, \dots, u_\ell = v)$ with

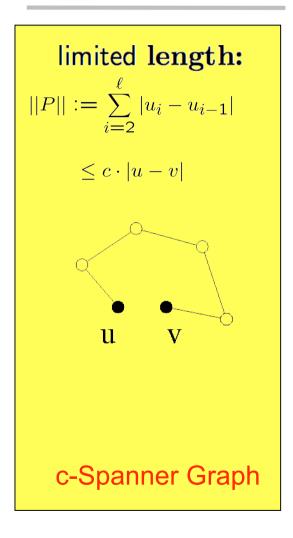




Related Work for Spanner Graphs

- Spanner Graphs introduced by Chew 1986
- Peleg & Schaffer first use in
 Distributed Comptuing 1989
- Applications to
 - Motion planning [Clarkson 1987]
 - Spanner Trees approximating MST [Yao 1982]
 - Used for FPAS for Traveling Salesman and related problems [Arora et al. 1998]
- Classic survey of spanners by Eppstein [2000]

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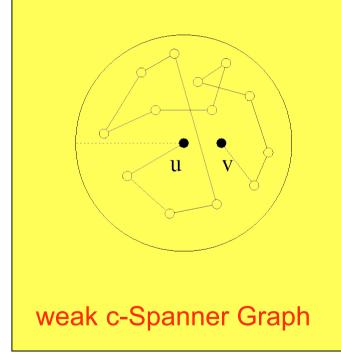
Related Work for Weak Spanners

- Applications in geometric searching and
- Constructions for spanners and weak spanners with arbitrary approximation ratio (stretch factor) by
 - Fischer, Meyer auf der Heide, Strothmann 1997
 - Fischer, Lukovszki, Ziegler 1998
- Optimization of routing time in wireless networks
 - Grünewald, Lukovski, S., Volbert 2002
 - Jia, Rajaraman, Scheideler 2003
- Constructions benefitting from locality properties for ad-hoc routing networks
 - Li, Wan, Wang 2001
 - Grünewald, Lukovski, S., Volbert 2002
 - Wang, Li 2002
 - Wang, Li, Wan, Frieder 2002

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-in a limited radius:

 $\max_{i=1,\dots,\ell} |u - u_i| \le c \cdot |u - v|$





Related Work for Power Spanners

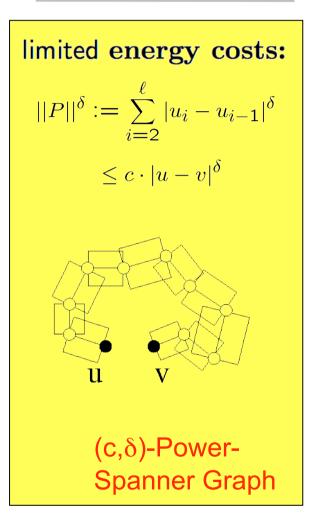
Power Spannners used in

- Grünewald, Lukovszki, S., Volbert 2002
- Li 2003
- Meyer auf der Heide, S, Volbert, Grünewald, 2004

> Special cases:

- δ = 0: Hop Spanners,
 - used by Alzoubi et al. 2003
- δ = 1: Power Spanner = Spanner
- δ = 2: Usual Power Spanner
 - (ad hoc networking)
- δ > 2: Energy consumption for messages in reality
 - Rappaport 1996 chooses δ up to 8







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Spanners versus Weak Spanners

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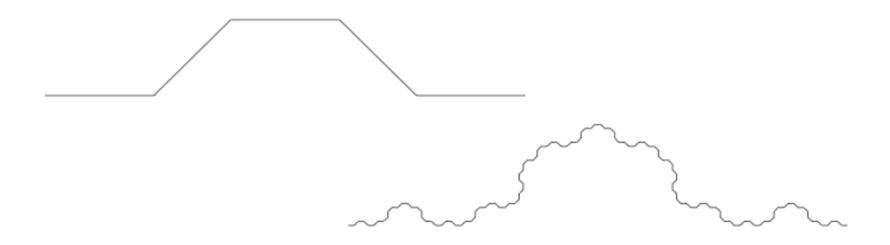
≻ Fact

- Every c-Spanner is also a c-Weak Spanner

Theorem

- There are Weak Spanner which are no Spanners

Proof Idea [Eppstein]: use fractal construction



Thank you

(and thanks go also to Holger Karl for providing some slides)



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